

Department of Mathematics, AGADU GE3(SB) Page-121
 allot availability as $\sum_{i=1}^n b_i = \sum_{j=1}^m a_j$ and the cost from the dummy origin to any destination is taken as zero.

Example 1 Find a basic feasible solution by North-West corner rule to the following transportation problem:

ORIGINS	DESTINATIONS					
	1	2	3	4	5	
1	2	11	10	3	7	4
2	1	4	7	2	1	8
3	3	9	4	8	12	9
	3	3	4	5	6	

Solution: Here $\sum_{i=1}^3 a_i = \sum_{j=1}^5 b_j = 21$, so it is a balanced TP.

Now we find a basic feasible solution by North-West corner rule, the steps are shown in the following tables:

	D ₁	D ₂	D ₃	D ₄	D ₅	a _i
O ₁	3		11	10	3	7
O ₂		1	4	7	2	1
O ₃			9	4	8	12
b _j	3	3	4	5	6	

→

	D ₂	D ₃	D ₄	D ₅	
O ₁	1		10	3	7
O ₂		1	7	2	1
O ₃		9	4	8	12
b _j	3	4	5	6	

→

	D ₂	D ₃	D ₄	D ₅	
O ₂	2		7	2	1
O ₃		9	4	8	12
b _j	3	4	5	6	

→

	D ₃	D ₄	D ₅	
O ₂	4		2	1
O ₃		9	8	12
b _j	4	5	6	

→

	D ₄	D ₅	
O ₂	3		1
O ₃		9	12
b _j	3	6	

→

	D ₅	
O ₃	6	12
b _j	6	

So, an initial basic feasible solution by North-West corner rule is $x_{11} = 3, x_{12} = 1, x_{22} = 2, x_{23} = 4, x_{24} = 2, x_{34} = 3$ and $x_{35} = 6$ and the

$$\begin{aligned} \text{Corresponding cost} &= 3 \times 2 + 1 \times 11 + 2 \times 4 + 4 \times 7 + 2 \times 2 + 3 \times 8 + 6 \times 12 \\ &= 6 + 11 + 8 + 28 + 4 + 24 + 72 \\ &= 153 \end{aligned}$$

Now, also we can write down the whole basic feasible solution in a table with this short-cut approach

	D ₁	D ₂	D ₃	D ₄	D ₅	a _i
O ₁	3 2	1 11				4X
O ₂		2 4	4 7	2 2		8/2
O ₃				3 8	6 12	9/6
	3	2	4	8	6	

Exercises: 1. Find an initial basic feasible solution by North West Corner method for the following transportation problem :

	D ₁	D ₂	D ₃	
O ₁	10	8	7	10
O ₂	6	9	8	5
	5	3	7	

2. Find an initial basic feasible solution by North West Corner rule for the following transportation problem

	D ₁	D ₂	D ₃	D ₄	
O ₁	3	7	2	1	11
O ₂	9	4	7	3	20
O ₃	10	2	8	3	35
	10	5	21	30	

Is the solution non-degenerate?

(2) Row minima method: In this method, we first consider the first row and find the minimum cost cell. Let $(1, k)$ cell be the cell in the first row with minimum cost. We allot in this cell the maximum allocation $x_{1k} = \min(a_1, b_k)$. If $a_1 < b_k$ then $x_{1k} = a_1$ and we cross out the first row and consider the remaining table and proceed in the same way. Again if $a_1 > b_k$ then $x_{1k} = b_k$ and we cross out the k th column and consider the remaining and consider the remaining table & proceed next in the same way. If $a_1 = b_k$, then both the first row and column k is cross and we proceed with the remaining table in the same way.

Example 2 Find an initial basic feasible solution by row minima method for the following transportation problem:

	D_1	D_2	D_3	D_4	D_5	
O_1	2	11	10	3	7	4
O_2	1	4	7	2	1	8
O_3	3	9	4	8	12	9
	3	3	4	5	6	

Solution: In the first row, minimum cost cell is $(1, 1)$

So, $x_{11} = \min(a_1, b_1) = \min(4, 3) = 3$. So the demand of the 1st destination is met and we delete the 1st column.

The remaining table is

	D_2	D_3	D_4	D_5	
O_1	11	10	3	7	1 (= 4-3)
O_2	4	7	2	1	8
O_3	9	4	8	12	9
	3	4	5	6	

In the first row of this table, the minimum cost cell is (1, 4) cell

$$\text{So, } x_{14} = \min(1, 5) = 1$$

Since the availability of the 1st row is exhausted, we delete the first row. The remaining table is

	D ₂	D ₃	D ₄	D ₅	
O ₂	4	7	2	1	8
O ₃	9	4	8	12	9
	3	4	4	6	

(=5-1)

Proceeding the same way, $x_{25} = \min(8, 6) = 6$. We delete destination D₅. The remaining table is

	D ₂	D ₃	D ₄	
O ₂	4	7	2	2 (=8-6)
O ₃	9	4	8	9
	3	4	4	

Now $x_{24} = \min(2, 4) = 2$. We delete row O₂

Next table is

	D ₂	D ₃	D ₄	
O ₃	9	4	8	9
	3	4	2	

(=4-2)

So, $x_{33} = \min(9, 4) = 4$. We delete destination D₃

Next table is

	D ₂	D ₄	
O ₃	9	8	5 (=9-4)
	3	2	

$x_{34} = \min(5, 2) = 2$. We delete destination D₄

Next table is

	D ₂	
O ₃	9	3 (=5-2)
	3	

So, $x_{32} = \min(3, 3) = 3$